Let X be a rv with support  

$$X = \begin{cases} \{x : p_{X}(x) > 0\} & \text{if } X & \text{is discrete with pmf } p_{X} \\ \{x : f_{X}(x) > 0\} & \text{if } X & \text{is continuous with pdf } f_{X} \end{cases}$$

Consider a function  $g: X \rightarrow Y$  taking values in X and returning values in some set Y, and define the  $vv \ Y$  to be Y = g(X).

Illustration:





Different points in X may map to the same point in Y.



$$I$$

$$V_{x_{1}}$$

$$V_{y_{1}}$$

Example: Let 
$$X \sim \text{Binamial}(n, p)$$
 and let  $Y = j(x)$ ,  
where  $g(x) = n - x$ . Recall that  
 $p_X(x) = \begin{cases} \binom{n}{x} p^X (1-p)^{n-x} & \text{if } x \in \{0, 1, ..., n\} \\ 0 & \text{otherwise.} \end{cases}$   
The transformation  $j(x) = n - x$  gives  $2f = \{0, 1, ..., n\}$ ,

The transformation 
$$j(x) = n - x$$
 gives  $2f = \{0, 1, ..., n\}$ ,  
and its inverse is defined by  
 $j^{-1}(y) = \{x \in \mathcal{X} : n - x = y\} = n - y$   
solve for  $x$ ,  
getting  $x = n - y$ 

So the part of Y bur 
$$y \in Q_{f}$$
 is given by  

$$p_{Y}(y) = \sum_{\substack{X \in p_{f}^{-1}(y) \\ X (x)}} p_{X}(x)$$

$$= p_{X}(n-y)$$

$$= \binom{n}{n-y} p^{n-y} \binom{n-(n-y)}{p} \binom{n}{(n-y)} = \frac{n!}{(n-y)!(n-(n-y)!)!} \binom{n}{(n-y)!(n-(n-y)!)!} \binom{n}{(n-y)!(y!)!} = \binom{n}{(y)}$$

$$= \binom{n}{y} \binom{(1-p)^{Y}}{p} p^{n-Y}.$$
So that Y ~ Binomial  $(n, 1-p)$ .  
Tradead, if X = st success, then  $Y = n - X = st$  follows in  
n independent Bernaulti trials with secuss probability  $p$ .  

$$\boxed{X \quad CONTTINUOUS}$$
Let X be continuous with pdf  $f_{Y}(x)$ . We first consider the edt of Y =  $g(X)$ , which is given by  
 $F_{Y}(y) = P_{Y}(Y = y)$ 

$$= P_{X}(y(X) = y)$$

$$= \int_{Y} (y(X) = y)$$

$$= \int_{Y} (x \in [x \in I : g(x) = y])$$

$$= \int_{Y} (x (x) dx .$$
Evaluation, the integral involves identifying the integration bounds over  $y \in Y_{1}$ . Then, if  $\frac{1}{2y} F_{Y}(y)$  is continuous over  $y \in Y_{1}$ .

$$f_{y}(y) = \int_{y} F_{y}(y) \quad \text{for } y \in \mathcal{Y}.$$

Obtaining the distribution of Y by concrypting out the above interaction to get the cell of Y is demaktines called the "Addition of USSTRESUTED FUNCTIONS."  
Example: Let 
$$X \sim \frac{1}{2} (x) = \frac{1}{2} x^2 \mathbf{1} (-i \in x \in i)$$
 and let  $Y = |x|$ .  
Find the pit  $\frac{1}{2} x^2 \mathbf{1} (-i \in x \in i)$  and let  $Y = |x|$ .  
Find the pit  $\frac{1}{2} x^2 \mathbf{1} (-i \in x \in i)$  and let  $Y = |x|$ .  
First find the odd  $F_{Y}(y)$  for  $y \in [0, 1]$ :  
 $F_{Y}(y) = F_{Y}(y \in y)$   
 $= \int_{X} (|x| \le y)$   
 $= \int_{X} (|x| \le y)$   
 $= \int_{Y} \frac{3}{2} x^2 dx$   
 $= \frac{1}{2} x^3 \Big|_{-y}^{Y}$   
 $= \frac{1}{2} x^3 \Big|_{-y}^{Y}$   
 $= \frac{y^3}{2}$ .  
So that  
 $F_{Y}(y) = \begin{cases} \frac{1}{2} , & y \ge 1 \\ 0 & y \le y \le 1 \end{cases}$  (Y is always  $\le 1$ )  
 $F_{Y}(y) = 3y^2 \mathbf{1} (0 \le y \le 1)$ .

Example: Let X~ Uniform (0,1) and let Y = -log X.  
Find the pdf fy of Y.  
We have Y = -log X, So 
$$g(x) = -log X$$
, giving  $Y = (0, 0)$ .  
First find the edf  $F_Y(y)$  for  $y \in (0, 0)$ :  
 $F_Y(y) = \int_{\{X \in \{0,1\}: -log X \le y\}} 1 \cdot dX = \int_{e^Y}^{1} 1 \cdot dX = 1 - e^Y$   
So  $x = e^Y$   
 $F_Y(y) = \begin{cases} 1 - e^{-Y}, & y = 0\\ 0, & y \le 0. \end{cases}$   
Taking  $\frac{d}{dy} F_Y(y)$  over  $y \ge 0$ , we get  
 $f_Y(y) = e^{-Y} \mathbb{1}(y \ge 0)$ .

The next theorem gives us the formula.

 $\frac{\text{Theorem}}{\text{theorem}}: \quad \text{Let } X \quad \text{be a continuous } rv \quad \text{with support } X \\ \text{and } pdf \quad f_X \quad \text{which is continuous on } X \\ \text{Then let } Y = j(X), \quad \text{such that } j: X \rightarrow Y \\ \text{is monotone and the derivative of } j: \\ \text{is continuovs on } Y \\ \text{is continuovs on } Y \\ \text{of } Y \quad \text{is } j: \\ \text{ven by } \\ f_Y(y) = \left\{ \begin{array}{c} f_X(\overline{g}'(y) \\ \overline{f}_Y \ \overline{f}'(y) \\ 0 \end{array} \right| \quad \frac{d}{f_Y} \ \overline{f}'(y) \\ 0 \end{array} \right\}, \quad y \in Y \\ \text{o therewise.} \end{array}$ 

Applying this theorem to get the density of Y is sometimes colled the "METHOD OF TRANSFORMATIONS"

**Proof:** Since 
$$j: I \neq ij$$
 is monotone over  $I$ , for any  $u, v \notin I$   
 $u \neq v \quad \langle z \geq g(u) \geq g(v)$  or  $u \geq v \quad \langle z \geq g(u) \leq g(v)$ .  
In this case the transformation is one-to-one, so that  
the inverse  $j^{-1}$  is single-valued, satisfying  
 $j^{-1}(v) = x \quad \langle z \geq y = g(x)$ .  
We consider the oft  $Y$  in the cases  
(i)  $g$  is monotone increasing:  
 $F_{Y}(v) = F_{Y}(Y \in y) = F_{X}(g(x) \leq y) = F_{X}(x \leq j^{-1}(v)) = F_{X}(g^{-1}(v))$   
(ii)  $g$  is monotone decreasing:  
 $F_{Y}(v) = F_{Y}(Y \in y) = F_{X}(g(x) \leq y) = F_{X}(x \geq j^{-1}(v)) = I - F_{X}(j^{-1}(v))$   
We have used here the castimity of  $X$ :  
 $F_{Y}(x) = F_{Y}(Y \in y) = I - F_{X}(x \leq j^{-1}(v)) = I - F_{X}(j^{-1}(v))$   
Then

Then

$$f_{Y}(y) = \frac{d}{dy} F_{Y}(y) = \begin{cases} \frac{d}{dy} F_{X}(j^{-1}(y)) & j \text{ increasing} \\ \frac{d}{dy} [1 - F_{X}(j^{-1}(y))] & j \text{ decreasing} \end{cases},$$
where by the chain rule we have
$$\frac{d}{dy} F_{X}(j^{-1}(y)) = f_{X}(j^{-1}(y)) \frac{d}{dy} j^{-1}(y) = f_{X}(j^{-1}(y)) \left| \frac{d}{dy} j^{-1}(y) \right|$$
and
$$\frac{d}{dy} F_{X}(j^{-1}(y)) = -f_{X}(j^{-1}(y)) \frac{d}{dy} j^{-1}(y) = f_{X}(j^{-1}(y)) \left| \frac{d}{dy} j^{-1}(y) \right|$$

$$\frac{d}{dy} \int [1 - F_{X}(j^{-1}(y))] = -f_{X}(j^{-1}(y)) \frac{d}{dy} j^{-1}(y) = f_{X}(j^{-1}(y)) \left| \frac{d}{dy} j^{-1}(y) \right|.$$

$$= 0, \text{ because } j \text{ decreasing}$$

$$= 2 j^{-1} \text{ decreasing}$$
Combining the two cases gives the result.

Example: Let 
$$X \sim f_X(x) = \frac{1}{\lambda} e^{-\frac{\gamma}{\lambda}} \mathbf{1}(x > 0)$$
. Find the post of  $Y = JX$ .  
We have  $g(x) = JX$ , which is monotone and gives  $Y = (0, \infty)$ .  
Also  $g^{-1}(y) = y^2$  (solve  $y = JX$  for  $x$ ), for which  
 $\frac{1}{\delta y} g^{-1}(y) = \frac{1}{\delta y} y^2 = 2y$ ,  
which is continuous over  $Y = (0, \infty)$ .

So applying (()) gives  

$$f_{y}(y) = \begin{cases} \frac{1}{\lambda} e^{-\frac{y^{2}}{\lambda}} |2y|, \quad y \ge 0 \\ 0, \quad y \le 0 \end{cases}$$
  
 $= (2/\lambda) y e^{-\frac{y^{2}}{\lambda}} \mathbf{1}(y \ge 0).$ 

Example: Let  $X \sim f_X(x) = \frac{1}{\Gamma(d)} p^{d_X} x^{-1} e^{-\frac{2\pi}{\beta}} \mathbf{1}(x \gg)$ . Find the pot of  $Y = x^{-1}$ . We have f(x) = 1/x is monotone on  $\mathcal{I} = (0, \infty)$  and gives  $\mathcal{Y} = (0, \infty)$ .

Also 
$$j'(y) = 1/y$$
 (solve  $y = 1/x$  for  $x$ ), for which

$$f_{y} = \frac{1}{y^{2}} = -\frac{1}{y^{2}} = -\frac{1}{y^{2}}$$

 $\frac{1}{4y} \frac{1}{y^{2}} \frac{1}{(y)} = -\frac{1}{y^{2}} = \frac{1}{y^{2}}$ which is continuous over  $\mathcal{Y} = (0, \infty)$ . So applying  $(\mathcal{O})$  sives

$$\begin{aligned} & = \frac{1}{\Gamma(4)\rho^{4}} \begin{pmatrix} \frac{1}{2} \end{pmatrix} \quad sives \\ & = \frac{1}{\Gamma(4)\rho^{4}} \begin{pmatrix} \frac{1}{2} \end{pmatrix}^{4-1} e^{-\frac{1}{2}\rho} \left| \frac{1}{2} \right|, \quad \gamma > 0 \\ & = \frac{1}{\Gamma(4)\rho^{4}} \begin{pmatrix} \frac{1}{2} \end{pmatrix}^{4+1} e^{-\frac{1}{2}\rho} \quad \mathbf{1}(\gamma > 0). \end{aligned}$$

q

Example: Let  $X \sim f_X(x) = e^{-x} \mathbf{1}(x \ge 0)$  and let  $Y = -\log X$ . Find the pdf and cdf of Y.

> We have  $y = f(x) = -\log x$  monotone on  $I = (0, \infty)$  and gives  $Y = (-\infty, \infty)$ . Also  $x = \overline{f}(y) = e^{-y}$ , for which  $\frac{d}{dy} \overline{f}(y) = -e^{-y}$ , which is continuous over  $Y = (-\infty, \infty)$ .

So applying (()) gives  $f_y(y) = e^{-y} | -e^{-y} | = e^{-(y+e^{-y})}$  for all  $y \in \mathbb{R}$ .

The edf of Y is given by

$$F_{y}(y) = \int_{-\infty}^{y} -(t+e^{-t}) dt = e^{-e^{-t}} = e^{-e^{-y}} \text{ for ell } y \in \mathbb{R}.$$

The distribution with this edt is called the Gumbel distribution. The Gumbel post looks like this:



## SQUARE TRANSFORMATION OF CONTINUOUS X

Let X be continuous with pdf fx with support over positive and negative numbers. Find the pdf of  $Y = X^2$ .

We have g(x) = x2, which is not monotone over X.

For yoo, we have

$$F_{Y}(v) = P_{Y}(Y \leq y)$$

$$= P_{X}(x^{2} \leq y)$$

$$= P_{X}(-i\overline{y} \leq x \leq J\overline{y})$$

$$= P_{X}(x \leq J\overline{y}) - P_{X}(x \leq -i\overline{y})$$

$$= F_{X}(J\overline{y}) - F_{X}(-i\overline{y}),$$

so that, by the chain rule,  $f_{y}(y) = \frac{d}{J_{y}} \left[ F_{x}(J_{y}) - F_{x}(-J_{y}) \right]$   $= \frac{d}{J_{x}} \left( I_{y} \right) \left( \frac{1}{2} \right) \frac{1}{J_{y}} - \frac{d}{J_{x}} \left( -J_{y} \right) \left( -\frac{1}{2} \right) \frac{1}{J_{y}} \cdot \frac{1}{J_{y}} \cdot \frac{1}{J_{y}} + \frac{1}{J_{y}} \left( -\frac{1}{2} \right) \frac{1}{J_{y}} \cdot \frac{1}{J_{y}} \cdot \frac{1}{J_{y}} \cdot \frac{1}{J_{y}} + \frac{1}{J_{y}} \left( -\frac{1}{2} \right) \frac{1}{J_{y}} \cdot \frac$ 

 $\begin{array}{rcl} & & & & \\ & & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & & \\ & & & \\ & & & \\ & & & \\ & & & & \\ & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & &$ 

This formula can be used to get the following very important result... <u>Application: Square of Normal(0,1) is Oi-gquare(1)</u>: Let X~ Normal(0,1). Find the pdf of  $Y=\chi^2$ . We have  $f_{\chi}(x) = \frac{1}{12\pi} e^{-\frac{\chi^2}{2}}$ , so for y>0, by  $(\Box)$ ,  $f_{\chi}(y) = \frac{1}{2} \frac{1}{12\pi} \left[ \frac{1}{12\pi} e^{-\frac{(\sqrt{y})^2}{2}} + \frac{1}{12\pi} e^{-\frac{(-\sqrt{y})^2}{2}} \right]$  $= \frac{1}{\sqrt{2\pi}} \frac{1}{\sqrt{y}} e^{-\frac{y}{2}}$ .

We tind we may write  $\begin{pmatrix} \Gamma(\frac{1}{2}) = J\overline{R}, \text{ see let } 8 \\ \text{trom STAT 511} \end{pmatrix}$  $t_{\gamma}(\gamma) = \frac{1}{F(\frac{1}{2})} \frac{(\frac{1}{2}-1)}{2} = \frac{-\frac{\gamma}{2}}{2} + \frac{(\frac{1}{2}-1)}{2} = \frac{-\frac{\gamma}{2}}{2} + \frac{1}{2} + \frac{(\gamma+1)}{2} = \frac{-\frac{\gamma}{2}}{2} + \frac{1}{2} + \frac{(\gamma+1)}{2} + \frac{\gamma}{2} = \frac{1}{2} + \frac{(\gamma+1)}{2} + \frac{\gamma}{2} + \frac{1}{2} + \frac{(\gamma+1)}{2} + \frac{\gamma}{2} + \frac{(\gamma+1)}{2} + \frac{(\gamma$ 

which we recognize as the polt of the Gamma (d=1/2, p=2) distribution, which is also called the Chi-square distribution with 1 degree of freedom.

## THE PROBABILITY INTEGRAL TRANSFORM

The probability integral transform turns any continuous rv into a Uniform (0,1) rv by passing it through its own cdf — the cdf is the "probability integral." We present it in the following theorem:

Theorem: Let X be a 
$$rV$$
 with a continuous of  $F_X$  and define  $Y = F_X(X)$ . Then Y has the uniform distribution on  $(0,1)$ .

Proof: \* Suppose 
$$F_X$$
 is monotone (we don't need this assumption, but it simplifies the proof).  
Then  $y = F_X(x) \iff F^{-1}(y) = x$ .  
Also,  $Y = (0,1)$ , since  $F_X(x) \in (0,1)$  for  $x \in X$ .

Illustration:



The cot of Y too 
$$y \in (0,1)$$
 is  
 $F_{y}(y) = P_{y}(Y = y) = P_{x}(F_{x}(x) = y) = P_{x}(x = F_{x}(y)) = y$ ,  
 $y_{y}^{m}$  quantile  
of  $y$ 

۶.

$$F_{y}(y) = \begin{cases} 2 & , & i \leq y \\ y & , & o \leq y \leq i \\ 0 & , & y \leq 0 \end{cases}$$

which is the Uniform (0,1) edf.

## RANDOM NUMBER GENERATION

We can apply the probability integral transform, in reverse, to generate random numbers. To generate X~Fx, do the following: (i) Generate U~ Un: form (0,1) (i) Set  $X = F_X^{-1}(U)$ , where  $F_X^{-1}(n) = \inf \{x: F_X(x) \ge n\}$ the probability integral transform in reverse. This is see that the edt of X=F\_X(u) is F\_X: WL  $P_{\mu}\left(F^{-1}(\nu)\leq x\right)=P_{\mu}\left(\nu\leq F_{x}(x)\right)=F_{x}(x).$ We need only to figure out how to generate numbers from the Uniform (0,1) distribution. <u>Se:</u> It is not easy to make a computer generate random numbers; it is in fact impossible, so we must content ourselves with pseudorandom numbers. ot nov, generator ot pseudorandom the "Mersenne Twisster," introduced by noto and Takuji Nishimura in The best, as of numbers is Makoto Matsumoto 1997. It generates pseudorandom Uniform (0,1) values: Throughout this paper, bold letters, such as  $\mathbf{x}$  and  $\mathbf{a}$ , denote word vectors, which are w-dimensional row vectors over the two-element field  $\mathbb{F}_2 = \{0, 1\}$ , identified with machine words of size w (with the least significant bit at the right).

machine words of size w (with the least significant bit at the right). The MT algorithm generates a sequence of word vectors, which are considered to be uniform pseudorandom integers between 0 and  $2^w - 1$ . Dividing by  $2^w - 1$ , we regard each word vector as a real number in [0,1]. The algorithm is based on the following linear recurrence

 $\mathbf{x}_{k+n} := \mathbf{x}_{k+m} \oplus (\mathbf{x}_k^u | \mathbf{x}_{k+1}^l) A, \quad (k = 0, 1, \cdots).$  (2.1)

Makoto Makumoto and Takuji Nishimura (1998). Mersenne Twister: a 623-dimensionally equidistributed uniform pseudo-random number zenerator. ACM Trans. Model. Comput. Simul. 8, 1,3-30.

Example: Generate a realization of  $X \sim Exponential(x)$ . We have  $F_{X}(x) = \begin{cases} 1 - e^{-\chi/\pi}, & x > 0 \\ 0, & x \leq 0, \end{cases}$ which has inverse (quantile function)  $F_{X}^{-1}(n) = -\lambda \log(1-n)$  (solve  $n = 1 - e^{-\chi/\pi}$  for x) for  $n \in (0, 1)$ . So generate  $U \sim United (0, 1)$  and set  $\chi = -\lambda \log(1-v)$ .

TRANSFORMATIONS WITH MOMENT GENERATING FUNCTIONS

In some situations, the easiest way to find  
the distribution of 
$$Y = g(X)$$
 is to compute  
 $M_y(t) = EEe^{tY} = Ee^{tg(X)}$ 

provided the expectation exists for all t in a neighborhood of zero. This strategy is most advantageous when j is a "shift-and-scele" transformation, that is, when j(x) = ax+b for some  $a, b \in \mathbb{R}$ .

This is due to the following result (In STAT 511). Theorem: For any constants a and b, the mgt of aX+b is

$$M_{X+b}(t) = e^{tb}M_{X}(at)$$

 $\frac{P_{res}f}{X}: \qquad M_{aX+b}(t) = \mathbb{E} \exp\left[t(aX+b)\right] = \mathbb{E} \exp\left[(ta)X\right] \exp\left[tb\right] = e^{tb}M_{X}(at)$ 

Finding the distribution of Y=z(X) by finding the met of Y is sometimes called the METHOD OF MOMENT GENERATING FUNCTIONS.

Example: Let 
$$X \sim Gramma (a, \beta)$$
 and let  $Y = X/\beta$ .  
We know that  $M_X(t) = (1 - \beta t)^{-d}$ , so  
 $M_Y(t) = M_{X/\beta}(t) = M_X(t/\beta) = (1 - \beta (t/\beta))^{-d} = (1 - t)^{-d}$ .  
And  $(1 - t)^{-d}$  is the myst of the Gramma  $(a, 1)$  dist.  
Example: Let  $X \sim Normal (p, \sigma^2)$  and let  $Z = X - p$ .  
We have  $M_X(t) = e^{pt + \sigma^2 t/2}$ , so  
 $M_Z(t) = M_X - p(t)$   
 $= e^{-pt} M_X(t/\beta)$   
 $= e^{-pt} m_X(t/\beta)$   
 $= e^{t/2} p(t) + \sigma^2(t)^2/2$   
 $= e^{t/2}$ ,  
which we recognize as the myst of the  
Normal  $(a, 1)$  distribution.